A Proximal Sampling Algorithm

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Introduction



Design and analysis of fast algorithms for sampling problems by leveraging tools from optimization.



(a) Optimization, $\min f(x)$



- A proximal sampling algorithm for nonconvex, semi-smooth and composite potentials
- Improved complexity to sample from a distribution $\varepsilon\text{-close}$ to the target distribution in KL, χ^2 and Rényi divergences
- Close interplay between sampling and optimization Proximal point framework

Problem: sample from $\nu(x) \propto \exp(-f(x))$

(A1) f is semi-smooth, i.e., there exist $\alpha_i \in [0,1]$ and $L_i > 0$, $i = 1, \ldots, n$, s.t.

$$||f'(u) - f'(v)|| \le \sum_{i=1}^{n} L_{\alpha_i} ||u - v||^{\alpha_i}, \quad \forall u, v \in \mathbb{R}^d$$

where f'(x) is in the Frechet subdiffernetial $\partial \phi(x)$; Examples: n = 11) $\alpha_1 = 1$, smooth, 2) $\alpha_1 = 0$, nonsmooth, 3) $0 < \alpha_1 < 1$, weakly smooth

(A2) ν satisfies log-Sobolev inequality (LSI) or Poincaré inequality (PI).

$$\mathsf{LSI:} \ H_{\nu}(\rho) \leq \frac{C_{LSI}}{2} J_{\rho}(\nu), \quad \mathsf{PI:} \ \mathbb{E}_{\nu}[(\psi - \mathbb{E}_{\nu}[\psi])^2] \leq C_{PI} \mathbb{E}_{\nu}[\|\nabla \psi\|^2]$$

Observations: ν is not necessarily log-concave, f is not necessarily convex.

Comparison

Source	Complexity	Assumption	Metric
Chewi et al.	$\tilde{\mathcal{O}}\left(\frac{C_{\mathrm{PI}}^{1+1/\alpha}L_{\alpha}^{2/\alpha}d^{2+1/\alpha}}{\varepsilon^{1/\alpha}}\right)$	weakly smooth $\alpha > 0$, Pl	Rényi
This work	$ ilde{\mathcal{O}}\left(C_{\mathrm{PI}}L_{lpha}^{2/(1+lpha)}d^2 ight)$	semi-smooth, Pl	Rényi

Table: Complexity bounds for sampling from non-convex semi-smooth potentials.

Source	Complexity	Assumption	Metric
Nguyen et al.	$\tilde{\mathcal{O}}\left(C_{\text{LSI}}^{1+\max\{\frac{1}{\alpha_i}\}}\left[\frac{n\max\{L_{\alpha_i}^2\}d}{\varepsilon}\right]^{\max\{\frac{1}{\alpha_i}\}}\right)$	weakly smooth $\alpha_i > 0$, LSI	KL
This work	$\tilde{\mathcal{O}}\left(C_{\mathrm{LSI}}\sum_{i=1}^{n}L^{2/(\alpha_{i}+1)}_{\alpha_{i}}d\right)$	semi-smooth, LSI	KL
This work	$\tilde{\mathcal{O}}\left(C_{\mathrm{PI}}\sum_{i=1}^{n}L^{2/(\alpha_{i}+1)}_{\alpha_{i}}d\right)$	semi-smooth, PI	Rényi

Table: Complexity bounds for sampling from non-convex composite potentials. $(\Box \mapsto (\Box)) \to (\Box) \mapsto (\Box) \mapsto (\Box) \to (\Box) \to$

Alternating Sampling Framework (ASF)

Joint distribution $\pi(x, y) \propto \exp[-f(x) - \frac{1}{2\eta} ||x - y||^2]$

Algorithm 1 ASF (Shen, Tian and Lee 2021)

1. Sample $y_k \sim \pi^{Y|X}(y \mid x_k) \propto \exp[-\frac{1}{2\eta} \|x_k - y\|^2]$ 2. Sample $x_{k+1} \sim \pi^{X|Y}(x \mid y_k) \propto \exp[-f(x) - \frac{1}{2\eta} \|x - y_k\|^2]$

Restricted Gaussian Oracle (RGO)

Given y, sample from

$$\pi^{X|Y}(\cdot|y) \propto \exp\left(-f(\cdot) - \frac{1}{2\eta} \|\cdot -y\|^2\right).$$

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Algorithm 2 ASF (Shen, Tian and Lee 2021)

- 1. Sample $y_k \sim \pi^{Y|X}(y \mid x_k) \propto \exp[-\frac{1}{2\eta} \|x_k y\|^2]$
- 2. Sample $x_{k+1} \sim \pi^{X|Y}(x \mid y_k) \propto \exp[-f(x) \frac{1}{2\eta} ||x y_k||^2]$

Restricted Gaussian Oracle (RGO)

Given y, sample from

$$\pi^{X|Y}(\cdot|y) \propto \exp\left(-f(\cdot) - \frac{1}{2\eta} \|\cdot -y\|^2\right).$$

Without an implementable and provable RGO, ASF is only conceptual.

Nontrivial

Proximal point framework: constructs a sequence of proximal problems

$$x_{k+1} \leftarrow \operatorname{prox}_{\eta f}(x_k) = \operatorname{argmin}_{x} \left\{ f(x) + \frac{1}{2\eta} \|x - x_k\|^2 \right\}$$
(*) (1)

E.g., Chambolle-Pock for saddle point, ADMM for distributed optimization

Algorithm 3 PPF

1.
$$y_k \leftarrow \underset{x}{\operatorname{argmin}} \frac{1}{2\eta} \|x - x_k\|^2 = x_k$$

2. $x_{k+1} \leftarrow \underset{x}{\operatorname{argmin}} \left\{ f_{y_k}^{\eta}(x) := f(x) + \frac{1}{2\eta} \|x - y_k\|^2 \right\}$

ASF for sampling \longleftrightarrow PPF for optimization

RGO in sampling \longleftrightarrow proximal mapping in optimization

Relaxed Proximal Bundle Method (L. and Monteiro 2021)

f is convex and Lipschitz continuous (nonsmooth, $\alpha_1 = 0$). Subgradient method.

Approximately solve (1) by the cutting-plane method (implementable)

$$z_{j} \leftarrow \mathsf{prox}_{\eta f_{j}}(x_{0}) = \min_{z} \left\{ f_{j}(z) + \frac{1}{2\eta} \|z - z_{0}\|^{2} \right\}, \quad z_{0} = x_{k}$$

where $f_{j}(z) = \max\{f(w) + \langle f'(w), z - w \rangle : w \in \{z_{0}, z_{1}, \dots, z_{j-1}\}\}$



Complexities: PPF $\mathcal{O}(\varepsilon^{-1}) \times \text{cutting-plane } \mathcal{O}(\varepsilon^{-1}) \implies \text{total } \mathcal{O}(\varepsilon^{-2}) \text{ optimal}$

implementable and provable

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RGO Implementation

RGO: given y, sample from $\exp(-f_y^{\eta}(x))$

Algorithm 4 RGO Rejection Sampling

- 1. Compute an approximate stationary point w of f_u^η
- 2. Generate sample $X \sim \exp(-h_1(x))$
- 3. Generate sample $U \sim \mathcal{U}[0, 1]$

4. If

$$U \le \frac{\exp(-f_y^{\eta}(X))}{\exp(-h_1(X))},$$

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10 / 25

then accept/return X; otherwise, reject X and go to step 2.

Proposal: $\exp(-h_1(x))$ where $h_1(x) \leq f_y^{\eta}(x)$

Rejection Sampling

 $X \sim \pi^{X|Y}(\cdot|y)$ and

$$\mathbb{P}(X \text{ is accepted}) = \mathbb{P}\left(U \le \frac{\exp(-f_y^{\eta}(X))}{\exp(-h_1(X))}\right)$$
$$= \frac{\int \exp(-f_y^{\eta}(x))dx}{\int \exp(-h_1(x))dx} \ge \frac{\int \exp(-h_2(x))dx}{\int \exp(-h_1(x))dx}$$
(2)

Want to find h_1 and h_2 such that

- i) sampling $\exp(-h_1(x))$ is easy,
- ii) $h_1(x) \le f_y^{\eta}(x) \le h_2(x)$,
- iii) (2) is bounded from below.

A Key Lemma

Consider n = 1, $\alpha \in [0, 1]$ and $L_{\alpha} > 0$

$$\|f'(u) - f'(v)\| \le L_{\alpha} \|u - v\|^{\alpha}, \quad \forall u, v \in \mathbb{R}^d;$$

Lemma

Assume f is L_{α} -semi-smooth, then for $\delta > 0$ and every $u, v \in \mathbb{R}^d$, we have

$$|f(u) - f(v) - \langle f'(v), u - v \rangle| \le \frac{M}{2} ||u - v||^2 + \frac{(1 - \alpha)\delta}{2}, \quad M = \frac{L_{\alpha}^{\frac{2}{\alpha+1}}}{[(\alpha + 1)\delta]^{\frac{1 - \alpha}{\alpha+1}}}.$$

Proof:

$$|f(u) - f(v) - \langle f'(v), u - v \rangle| \le \frac{L_{\alpha}}{\alpha + 1} ||u - v||^{\alpha + 1}$$

Young's inequality $ab \leq \frac{a^p}{p} + \frac{b^q}{q}$, $\frac{1}{p} + \frac{1}{q} = 1$ with

$$a = \frac{L_{\alpha}}{(\alpha+1)\delta^{\frac{1-\alpha}{2}}} \|u-v\|^{\alpha+1}, \quad b = \delta^{\frac{1-\alpha}{2}}, \quad p = \frac{2}{\alpha+1}, \quad q = \frac{2}{1-\alpha}.$$

12 / 25

Definition

A stationary point
$$w^*$$
 of f_y^{η} is such that $f'(w^*) + \frac{1}{n}(w^* - y) = 0$.

Definition

An approximate stationary point w of f_y^{η} is s.t. $\|f'(w) + \frac{1}{\eta}(w-y)\| \leq \sqrt{Md}$.

$$h_1(x) := f(w) + \langle f'(w), x - w \rangle - \frac{M}{2} ||x - w||^2 + \frac{1}{2\eta} ||x - y||^2 - \frac{(1 - \alpha)\delta}{2},$$

$$h_2(x) := f(w^*) + \langle f'(w^*), x - w^* \rangle + \frac{M}{2} ||x - w^*||^2 + \frac{1}{2\eta} ||x - y||^2 + \frac{(1 - \alpha)\delta}{2}.$$

Answers: i) sampling $\exp(-h_1(x))$ is easy; ii) verify $h_1(x) \le f_y^\eta(x) \le h_2(x)$ by the key lemma. Q1. Rejection sampling complexity

$$[\mathbb{P}(X \text{ is accepted})]^{-1} \leq \frac{\int \exp(-h_1(x)) dx}{\int \exp(-h_2(x)) dx}$$

Q2. Optimization complexity to find an approx. stat. pt. w s.t.

$$\left\|f'(w) + \frac{1}{\eta}(w-y)\right\| \le \sqrt{Md}$$

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Proposition

Assume

$$\eta \leq \frac{1}{Md} = \frac{\left[(\alpha+1)\delta\right]^{\frac{1-\alpha}{\alpha+1}}}{L_{\alpha}^{\frac{2}{\alpha+1}}d},$$

then the expected number of rejection steps in Algorithm 4 is at most $\exp\left(\frac{3(1-\alpha)\delta}{2}+3\right)$.

Intuition: if η is small enough, h_1 and h_2 are convex quadratic functions, so

$$\frac{\int \exp(-h_1(x))dx}{\int \exp(-h_2(x))dx} \approx \left(\frac{1+\eta M}{1-\eta M}\right)^{d/2} \le (1+4\eta M)^{d/2} \le \left(1+\frac{4}{d}\right)^{d/2} \le e^2.$$

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Answer to Q2 – Optimization complexity

Lemma

Let
$$f_y^\eta := f + \frac{1}{2\eta} \|\cdot -y\|^2$$
 and $(f_y^\eta)' := f' + \frac{1}{\eta}(\cdot -y)$, then for every $u, v \in \mathbb{R}^d$,

$$\frac{1}{2}\left(\frac{1}{\eta} - M\right) \|u - v\|^2 - \frac{(1 - \alpha)\delta}{2} \le f_y^{\eta}(u) - f_y^{\eta}(v) - \langle (f_y^{\eta})'(v), u - v \rangle \\ \le \frac{1}{2}\left(\frac{1}{\eta} + M\right) \|u - v\|^2 + \frac{(1 - \alpha)\delta}{2}.$$

 f_y^η is nearly $(\eta^{-1}-M)\text{-strongly convex and }(\eta^{-1}+M)\text{-smooth}$

Proposition

Assume $\eta \leq \frac{1}{Md}$, then the iteration-complexity to find the approx. stat. pt. w s.t. $\left\|f'(w) + \frac{1}{\eta}(w-y)\right\| \leq \sqrt{Md}$ by Nesterov acceleration is $\tilde{\mathcal{O}}(1)$.

$$\mu = \frac{1}{\eta} - M \approx M(d-1), \quad L = \frac{1}{\eta} + M \approx M(d+1), \quad \sqrt{L/\mu} \approx 1$$

RGO and ASF Complexities

Putting previous results together, we can implement RGO with $\tilde{\mathcal{O}}(1)$ subgradients of f and $\mathcal{O}(1)$ samples from Gaussian distribution in expectation.

Other ingredients for total complexity: Convergence rate analysis of ASF

Theorem (Chen, Chewi, Salim and Wibisono 2022)

If $\nu \propto \exp(-f)$ satisfies LSI with $C_{LSI} > 0$, then x_k of ASF $\sim \rho_k$, which satisfies

$$H_{\nu}(\rho_k) \leq \frac{H_{\nu}(\rho_0)}{\left(1 + \frac{\eta}{C_{LSI}}\right)^{2k}}.$$

Theorem (Chen, Chewi, Salim and Wibisono 2022)

If $\nu \propto \exp(-f)$ satisfies PI with $C_{\rm PI}>0$, then x_k of ASF $\sim
ho_k$, which satisfies

$$\chi_{\nu}^{2}(\rho_{k}) \leq \frac{\chi_{\nu}^{2}(\rho_{0})}{\left(1 + \frac{\eta}{C_{\mathrm{PI}}}\right)^{2k}}.$$

Theorem

Suppose f is L_{α} -semi-smooth and ν satisfies PI. With $\eta \asymp 1/(L_{\alpha}^{\frac{2}{\alpha+1}}d)$, then ASF with RGO by rejection has complexity bound

$$\tilde{\mathcal{O}}\left(C_{\mathrm{PI}}L_{\alpha}^{\frac{2}{\alpha+1}}d\right)$$

to achieve ε error to ν in terms of χ^2 divergence. Each iteration queries $\tilde{\mathcal{O}}(1)$ subgradients of f and generates $\mathcal{O}(1)$ samples in expectation from Gaussian distribution.

$$\|f'(u) - f'(v)\| \le \sum_{i=1}^n L_{\alpha_i} \|u - v\|^{\alpha_i}, \quad \forall u, v \in \mathbb{R}^d;$$

Theorem

Suppose f is semi-smooth and ν satisfies LSI. With $\eta \asymp \left[\sum_{i=1}^{n} L_{\alpha_i}^{\frac{2}{\alpha_i+1}}d\right]^{-1}$, then ASE with PCO by rejection has complexity bound

ASF with RGO by rejection has complexity bound

$$\tilde{\mathcal{O}}\left(C_{\mathrm{LSI}}\sum_{i=1}^{n}L_{\alpha_{i}}^{\frac{2}{\alpha_{i}+1}}d\right)$$

to achieve ε error to ν in terms of χ^2 divergence. Each iteration queries $\tilde{\mathcal{O}}(1)$ subgradients of f and generates $\mathcal{O}(1)$ samples in expectation from Gaussian distribution.

Theorem

Suppose f is semi-smooth and ν satisfies PI. With $\eta \asymp \left[\sum_{i=1}^{n} L_{\alpha_i}^{\frac{2}{\alpha_i+1}} d\right]^{-1}$, then ASF with RGO by rejection has complexity bound

$$\tilde{\mathcal{O}}\left(C_{\mathrm{PI}}\sum_{i=1}^{n}L_{\alpha_{i}}^{\frac{2}{\alpha_{i}+1}}d\right)$$

to achieve ε error to ν in terms of χ^2 divergence. Each iteration queries $\hat{\mathcal{O}}(1)$ subgradients of f and generates $\mathcal{O}(1)$ samples in expectation from Gaussian distribution.

Interpretation of Unadjusted Langevin Algorithm (ULA)

Algorithm 5 ASF

- 1. Sample $y_k \sim \pi^{Y|X}(y \mid x_k) \propto \exp[-\frac{1}{2\eta} \|x_k y\|^2]$
- 2. Sample $x_{k+1} \sim \pi^{X|Y}(x \mid y_k) \propto \exp[-f(x) \frac{1}{2\eta} ||x y_k||^2]$

Algorithm 6 ULA

- 1. Sample $y_k \sim \pi^{Y|X}(y \mid x_k) \propto \exp[-rac{1}{2\eta} \|x_k y\|^2]$
- 2. Sample $x_{k+1} \sim e^{-\langle \nabla f(y_k), x y_k \rangle \frac{1}{2\eta} \|x y_k\|^2} \propto e^{-\frac{1}{2\eta} \|x (y_k \eta \nabla f(y_k))\|^2}$

$$\begin{aligned} x_{k+1} &= y_k - \eta \nabla f(y_k) + \sqrt{\eta} z_k, \quad z_k \sim N(0, I), \\ y_{k+1} &= x_{k+1} + \sqrt{\eta} z'_k, \quad z'_k \sim N(0, I). \end{aligned}$$

 $\implies y_{k+1} = y_k - \eta \nabla f(y_k) + \sqrt{\eta} (z_k + z'_k) = y_k - \eta \nabla f(y_k) + \sqrt{2\eta} z, \quad z \sim N(0, I)$

ULA can be viewed as ASF with RGO implemented without rejection

$$h_1(x) = f(y_k) + \langle \nabla f(y_k), x - y_k \rangle + \frac{1}{2\eta} \|x - y_k\|^2 \le f(x) + \frac{1}{2\eta} \|x - y_k\|^2 = f_{y_k}^{\eta}(x)$$

- A proximal sampling algorithm for $\nu \propto \exp(-f)$.
 - f nonconvex, semi-smooth, composite. ν satisfies either LSI or PI.

• Total complexity
$$\tilde{\mathcal{O}}\left(C\sum_{i=1}^{n}L_{\alpha_{i}}^{\frac{2}{\alpha_{i}+1}}d\right)$$
 where $C = C_{\text{LSI}}$ or $C = C_{\text{PI}}$.

Each iteration takes $\hat{\mathcal{O}}(1)$ subgradients of f and $\mathcal{O}(1)$ samples from Gaussian.

Inspired by proximal point framework and proximal mapping.
 Leverage tools from optimization to design and analyze sampling algorithms.
 E.g., acceleration in sampling for weakly smooth potentials.

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Thank you!

Definition (Frechet ε -subdiffernetial)

Let $\phi:\mathbb{R}^n\to\mathbb{R}\cup\{\infty\}$ be a proper closed function, then the Frechet ε subdiffernetial is defined as

$$\tilde{\partial}_{\varepsilon}\phi(x) = \left\{ v \in \mathbb{R}^n : \liminf_{y \to x} \frac{\phi(y) - \phi(x) - \langle v, y - x \rangle}{\|y - x\|} \ge -\varepsilon \right\}$$

When $\varepsilon = 0$, we denote $\tilde{\partial}_0 \phi(x)$ simply by $\tilde{\partial} \phi(x)$.

Lemma

If $\phi: \mathbb{R}^n \to \mathbb{R} \cup \{\infty\}$ is an *m*-weakly convex function, then for any $x, c \in \mathbb{R}^n$, we have

$$\tilde{\partial}\phi(x) = \partial\left(\phi_m(\cdot;c)\right)(x) - m(x-c) \tag{3}$$

where

$$\phi_m(\cdot; c) := \phi(\cdot) + \frac{m}{2} \| \cdot -c \|^2.$$

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